

## Final Exam Answers: Economics 101

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### 1. Normal Form Games (note that a complete answer must include a drawing of the socially feasible sets)

a)

	L( $p$ )	R( $1-p$ )
U( $q$ )	2*,5*	-1,1
D( $1-q$ )	1,-1	5*,2*

Two pure strategy equilibria as marked. Mixed for player 2  $2p - (1-p) = p + 5(1-p)$  so  $p=6/7$ ; for player 1  $5q - (1-q) = q + 2(1-q)$  so  $q=3/7$ . Pure strategy equilibria are Pareto Efficient. The mixed equilibrium has payoffs of  $(11/7, 11/7)$  is not. No weakly dominated strategies. Pure strategy maxmin for both players is 1; pure strategy minmax for both players is 2.

b)

	L	R
U	5,5	-1,8*
D	8*,-1	1*,1*

Unique Nash equilibrium (U,L are both strictly dominated). No mixed equilibria due to dominance. Nash equilibrium is not Pareto efficient. Pure maxmin and maxmin is 1 for both players.

c)

	L	R
U	-1*,3	-3,5*
D	-3,5*	-1*,3

No pure strategy equilibrium. Unique Pareto efficient mixed equilibrium where both players mix 50-50. No weakly dominated strategies. Note that the socially feasible set is one-dimensional. Pure strategy maxmin for player 1 is  $-3$ , for player 2 is  $3$ ; pure strategy minmax for player 1 is  $-1$ , for player 2 is  $5$ .

### 2. Long Run versus Long Run

	L	R
U	3,3	0,5
D	5,0	1,1

Use the grim strategies: U(or L) as long as UL in every past period, otherwise DR (the static Nash equilibrium). In equilibrium you get 3. If you deviate you get at most  $(1-\delta)5 + \delta 1 \leq 3$  or  $2 \leq 4\delta$ , so this is an equilibrium for  $\delta \geq 1/2$ .

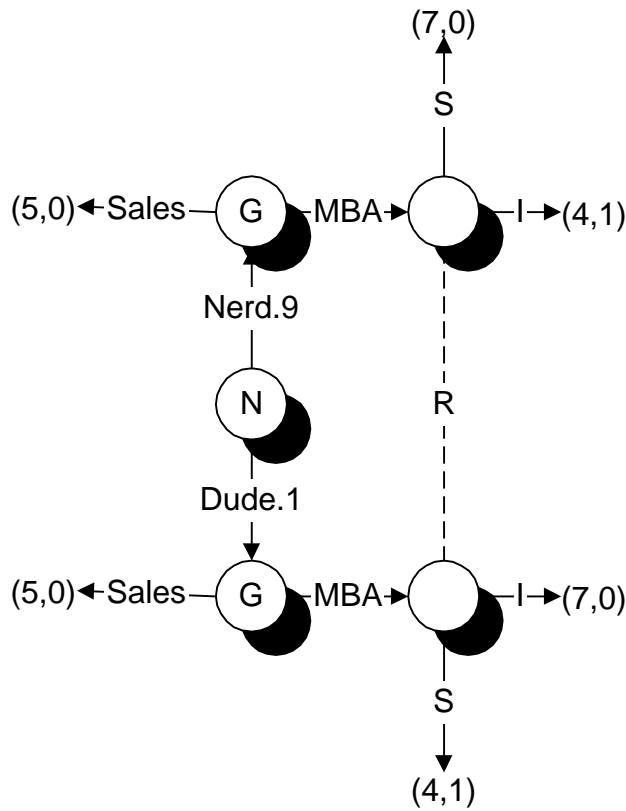
### 3. Long Run versus Short Run

	L	R
U	2,1*	0,0
D	11*,0	1*,3*

The unique Nash equilibrium is DR; the Stackelberg equilibrium is UL. Strategies for which lead to playing UL are UL if always UL in the past and DR if ever a deviation. Alternatively, players may base their strategies on past play of the LR player only: LR: U if U in the past and D if ever a deviation by LR and SR: L if U in the past and R if ever a deviation of the LR player.

These are optimal for the short-run player because it is in his best-response correspondence. For the long run player it must be that  $2 \geq (1-\delta)11 + \delta 1$  or  $\delta \geq 9/10$ .

#### 4. Screening



Nerd/Dude	$S(q)$	$I(1-q)$
SS	5,0*	5,0*
SM	4.9,0.1*	5.2*,0
MS	6.8*,0	4.1,0.9*
MM	6.7,0.1	4.3,0.9*

No pure equilibria. Observe that if player 1 randomizes with weight .9 on SM and .1 on MS, he gets 5.09 regardless of how player 2 plays. So SS is strictly dominated and will not be played. Next observe that for player 2 to mix, player 1 must put some weight on SM. Suppose that 1 is indifferent between SM and MS. Then  $4.9q + 5.2(1-q) = 6.8q + 4.1(1-q)$ . This gives  $q = 11/30$ , and the expected utility is 5.09. On the other hand, the expected utility from MM is 5.18. So next we try to make player 1

indifferent between SM and MM. Then  $4.9q + 5.2(1 - q) = 6.7q + 4.3(1 - q)$ , or  $q = 1/3$ , with an expected utility of 5.1. In this case the utility from MS is only 5. So we conclude that  $q = 1/3$ , and that player 1 is indifferent between SM and MM, and will not play SS or MS. Finally, to make player 2 indifferent, player 1 must choose the probability  $p$  of MM so that  $0.9p = 0.1$ , or  $p = 1/9$ .

What then is the probability of nerd|mba?

$p(n|m) = \frac{p(m|n)p(n)}{p(m)}$  The probability of  $n$  is .9. The probability of  $m$  is equal to  $1/9$  (the

probability of MM) plus  $8/9 \times .1$  (the probability of SM times probability of dude). The probability of  $m|n$  is  $1/9$ , since nerds stay out when SM is played. So

$$p(n|m) = \frac{(1/9) \cdot .9}{1/9 + .1 \times 8/9} = \frac{.9}{1 + .8} = 1/2$$

## 5. Price Discrimination

a)

$$(5 - p^H)x^H \geq (5 - p^L)x^L \text{ or } 5(x^H - x^L) \geq p^H x^H - p^L x^L$$

$$(3 - p^L)x^L \geq (3 - p^H)x^H \text{ or } 3(x^H - x^L) \leq p^H x^H - p^L x^L$$

Important observation: these two inequalities can be satisfied only if  $x^H \geq x^L$ . This in turn shows that  $p^H x^H \geq p^L x^L$ .

b)

$$(5 - p^H)x^H \geq 0 \text{ or } 5 \geq p^H$$

$$(3 - p^L)x^L \geq 0 \text{ or } 3 \geq p^L$$

c)

$$U = 5p^H x^H + 5p^L x^L$$

case 1:  $x^H = x^L$ ; then from a) we see that  $p^H = p^L$ , so  $U = p^L x^L$ . From b) we see that  $p^L \leq 3$ , so utility will be a maximum when  $p^L = 3$  and  $x^L = 2$ , yielding a utility of 6.

Case 2: Since from a)  $x^H \geq x^L$  the other case is  $x^H = 2$ ,  $x^L = 1$ . The constraints are

$5 \geq 2p^H - p^L, 2p^H - p^L \geq 3, 5 \geq p^H, 3 \geq p^L$ , utility is  $U = p^H + 5p^L$ .

Rewrite constraints  $(5 + p^L)/2 \geq p^H, 2p^H - 3 \geq p^L$  so

$$p^H \leq \max\{(5 + p^L)/2, 5\}$$

$$p^L \leq \max\{2p^H - 3, 3\}$$

Case 1a)  $p^H = 5$  then from second constraint  $p^L = 3$ , which means that  $p^H \leq 4$  so this case is not possible.

Case 1b)  $p^H = (5 + p^L)/2$  then  $p^L \leq (5 + p^L) - 3 = p^L + 2$ , which doesn't bind, so  $p^L = 3$ . Then  $p^H = 4$ , which satisfies  $p^H \leq \max\{(5 + p^L)/2, 5\}$ .

Utility is then  $U = p^H + 5p^L = 4 + 15 = 5.5$

So we should sell at the fixed price of 3 and not try to price discriminate.